

# EXHIBIT A

**United States Patent** [19]  
**Snodgrass et al.**

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 [45] **Date of Patent:** **Mar. 19, 1996**

- [54] **DATA COMMUNICATION METHOD USING IDENTIFICATION PROTOCOL**
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- [73] Assignee: **Micron Technology, Inc.**, Boise, Id.
- [21] Appl. No.: **990,915**
- [22] Filed: **Dec. 15, 1992**
- [51] Int. Cl.<sup>6</sup> ..... **G01S 13/76**
- [52] U.S. Cl. .... **342/42; 342/44; 342/50; 342/51**
- [58] **Field of Search** ..... **342/42, 44, 45, 342/50, 51**

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[57] **ABSTRACT**

A protocol is used to coordinate the use of a common communication medium by one or more interrogating commander stations and an unknown plurality of responding responder stations. Each commander station and each responder station is equipped to broadcast messages and to check for error in received messages. When more than one station attempts to broadcast simultaneously, an erroneous message is received and communication is interrupted. To establish uninterrupted communication, a commander station broadcasts a command causing each responder station of a potentially large first number of responder stations to each select a random number from a known range and retain it as its arbitration number. After receipt of such a command, each addressed responder station transmits a response message containing its arbitration number. Zero, one, or several responses may occur simultaneously. By broadcasting requests for identification to various subsets of the full range of arbitration numbers and checking for an immediate error-free response, a commander station can determine the arbitration number of every responder station capable of communicating at the time. Consequently, a commander station can conduct subsequent uninterrupted communication with each responder station, for example by addressing only one responder station. Responder stations of this invention require minimal logic and circuitry to respond to multiple commander stations.

**19 Claims, 10 Drawing Sheets**

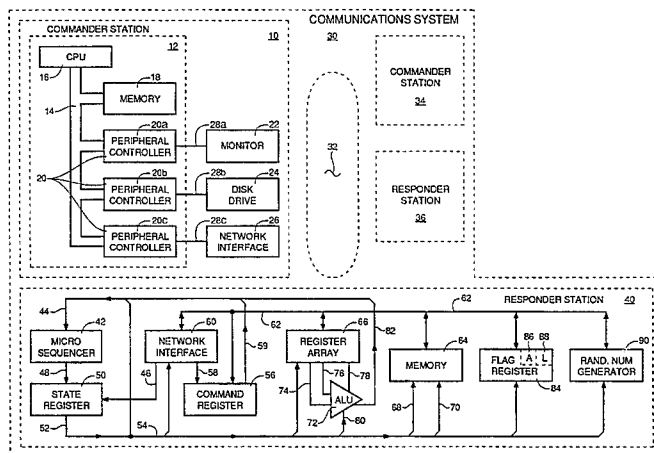
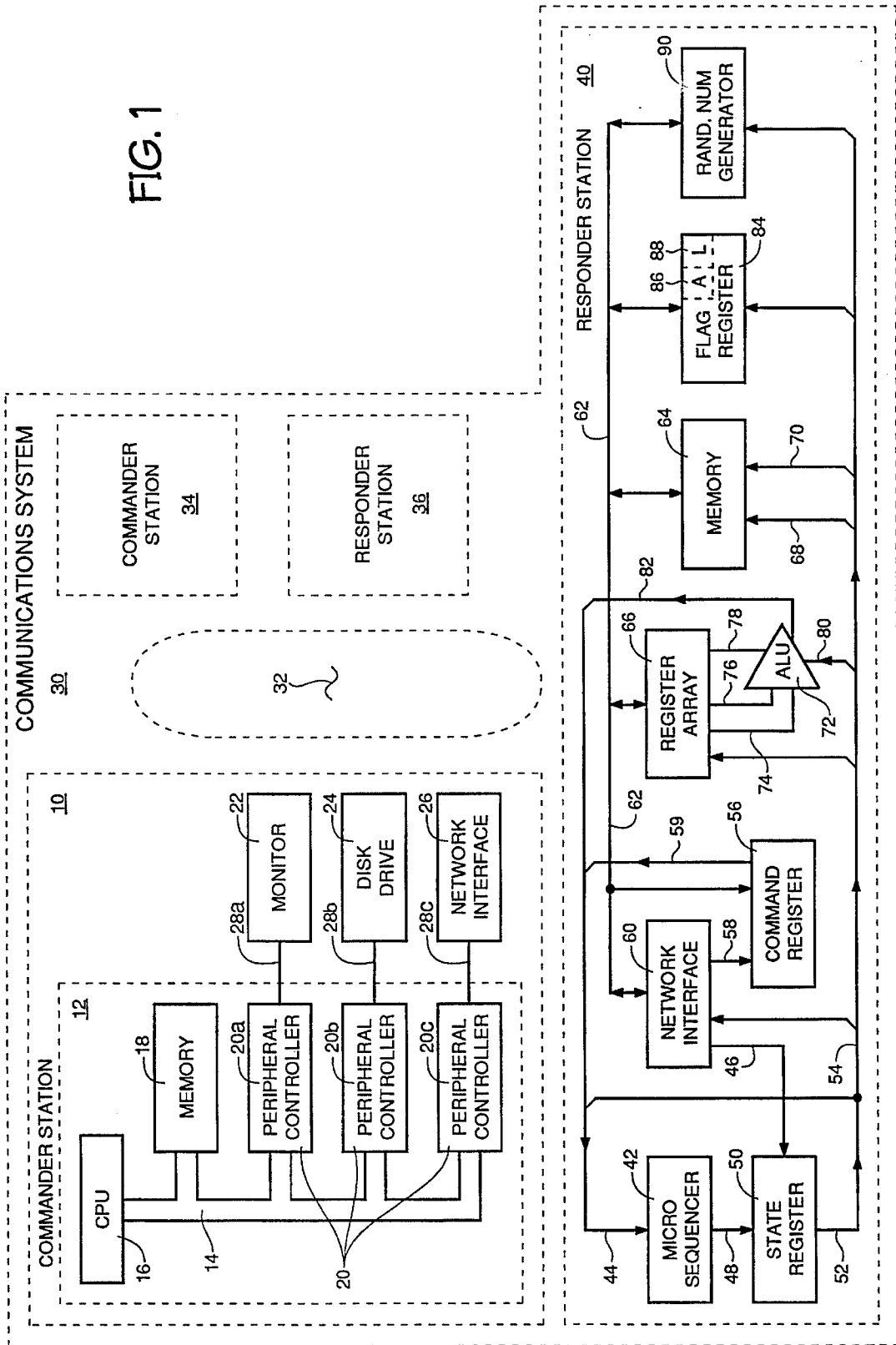


FIG. 1



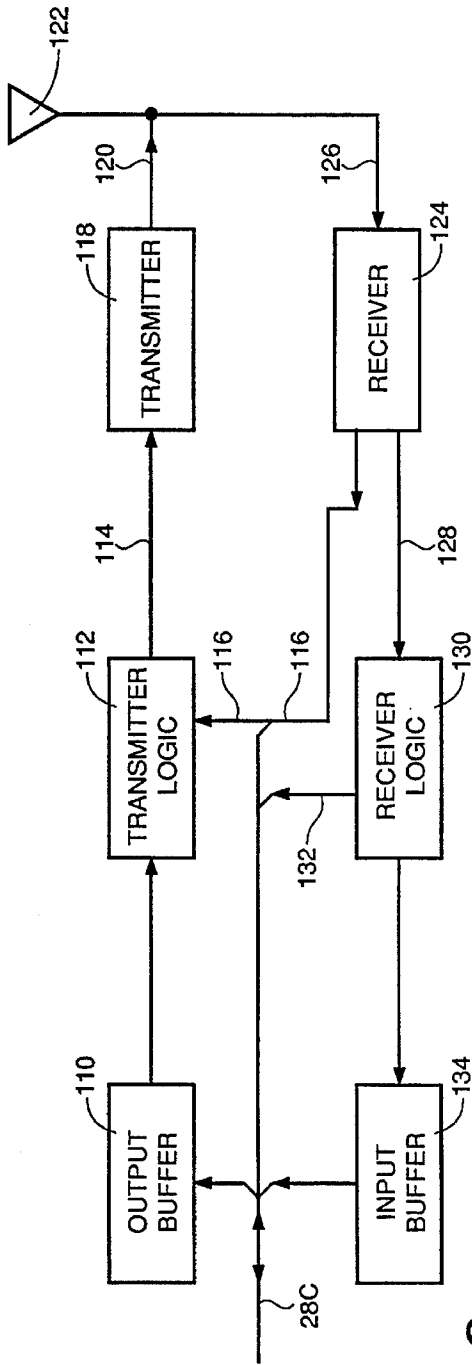


FIG. 2

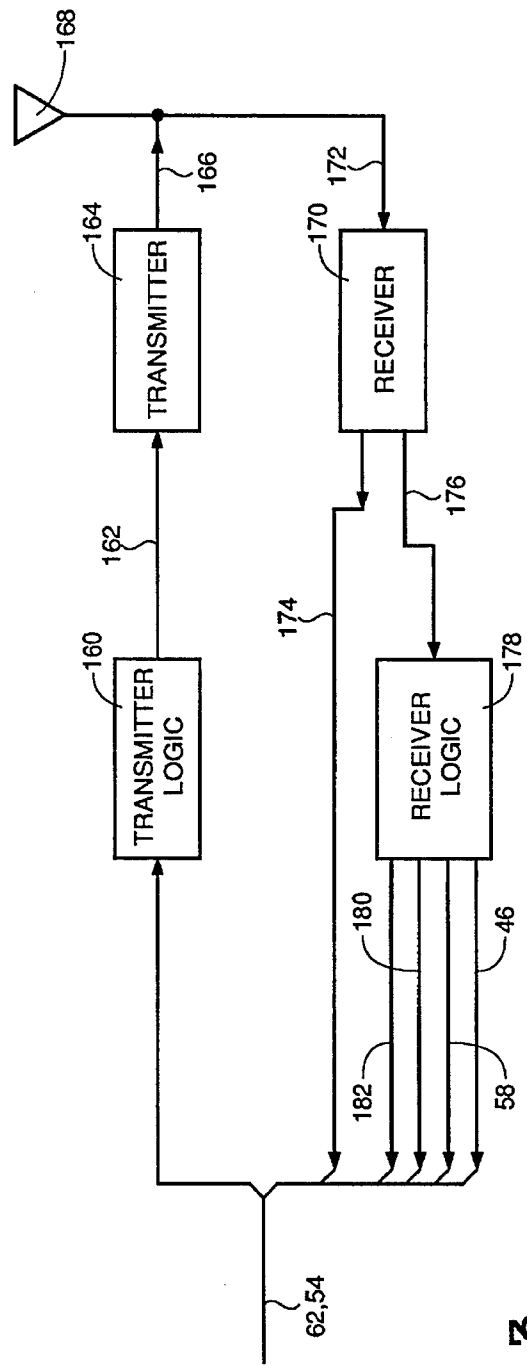


FIG. 3



FIG. 4

COMMANDS			
OPCODE (HEXADECIMAL)	MNEMONIC	DESCRIPTION	FORMAT
08	ID	IDENTIFY	142
A9	IDG	IDENTIFY AND GENERATE ARBITRATION NUMBER	142
EA	IDC	IDENTIFY AND CLEAR LOCKED BIT 88	142
4B	IDCG	IDENTIFY, CLEAR LOCKED BIT 88, AND GENERATE ARBITRATION NUMBER	142
8E	RD	READ DATA AND SET LOCKED BIT 88	144
2F	WD	WRITE DATA AND SET LOCKED BIT 88	146

FIG. 5

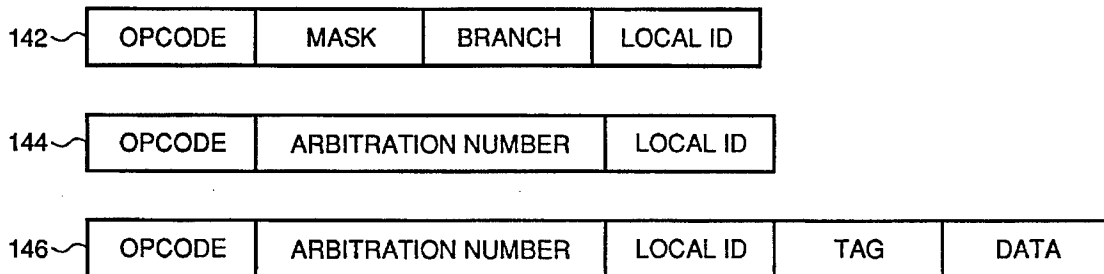


FIG. 6

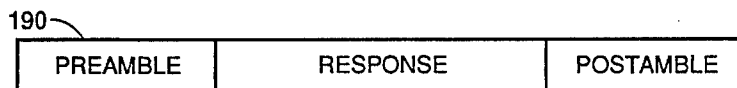


FIG. 7

RESPONSES		
MNEMONIC	DESCRIPTION	FORMAT
IDR	RESPONSE TO ID, IDG, IDC, AND IDCG COMMANDS	192
RDR	RESPONSE TO RD COMMAND	194
WDR	RESPONSE TO WR COMMAND	196

FIG. 8

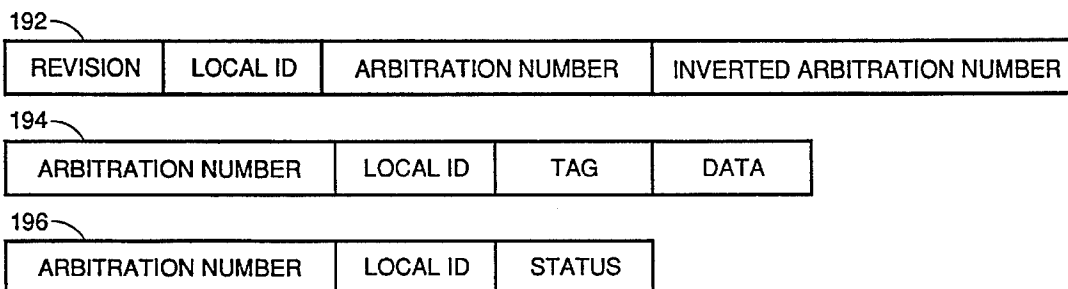


FIG. 9

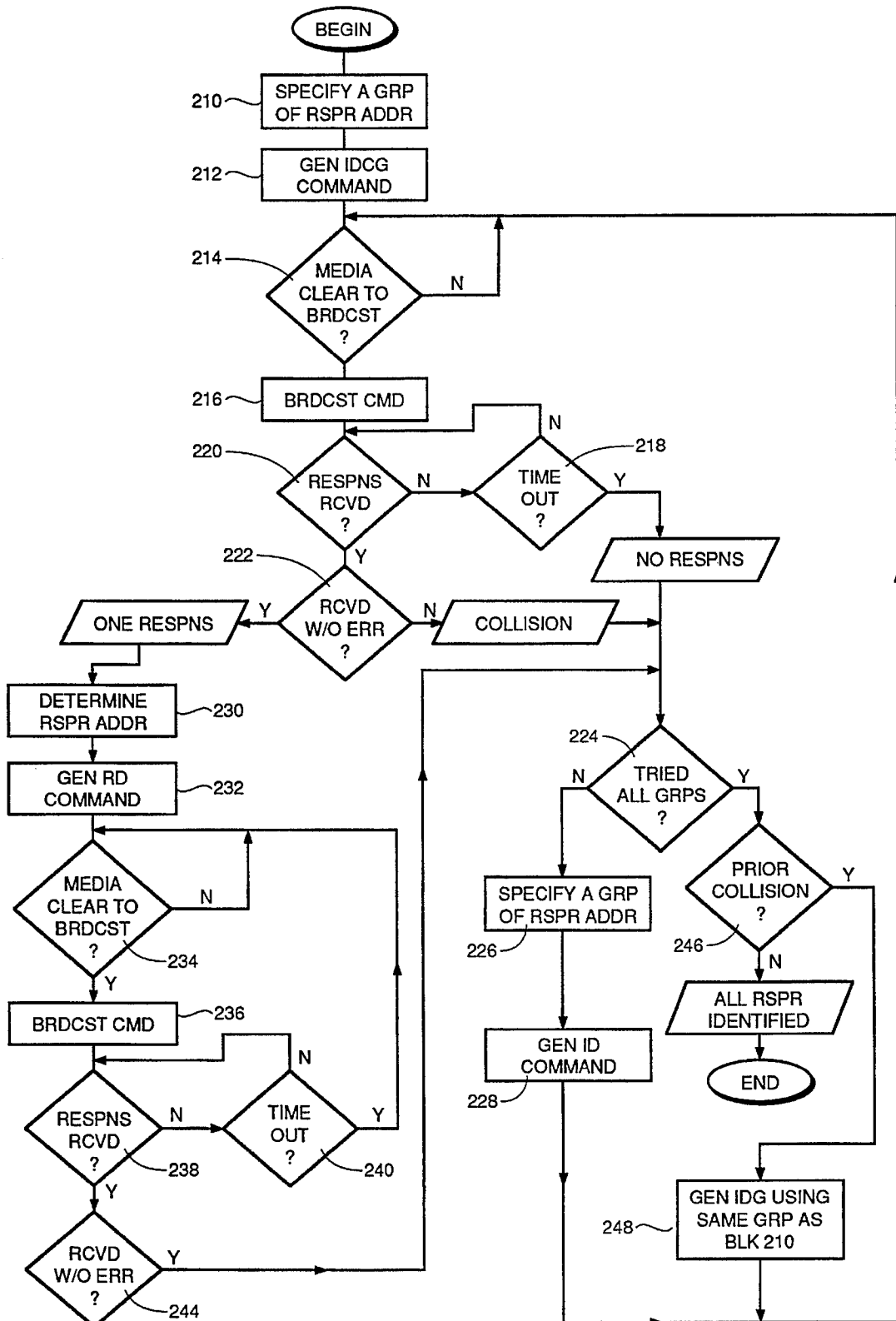


FIG. 10

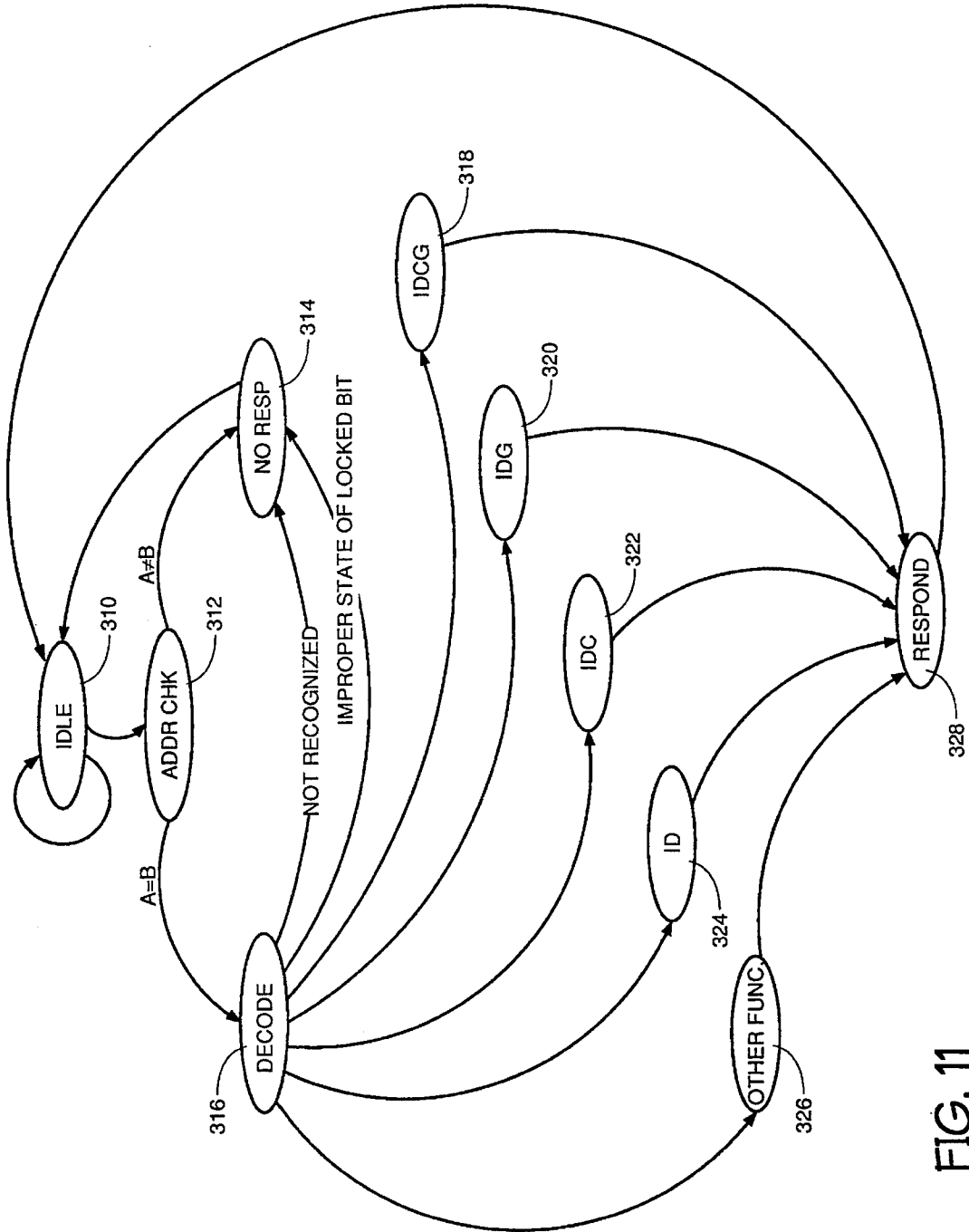


FIG. 11



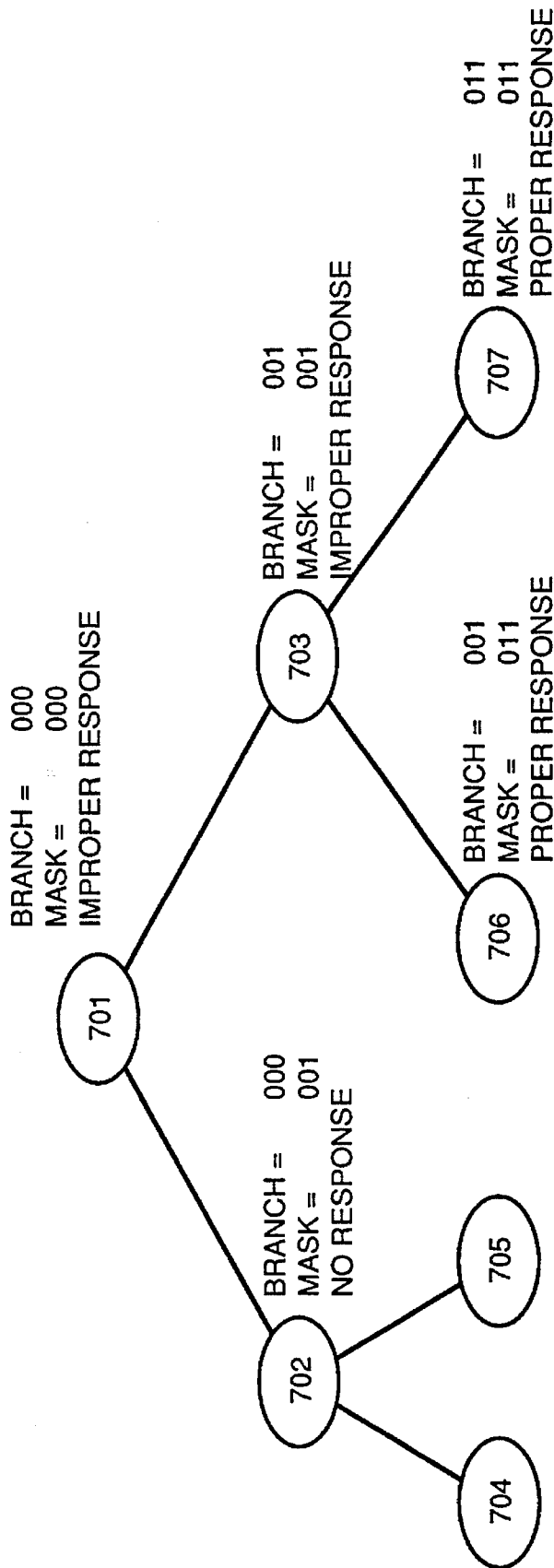


FIG. 12

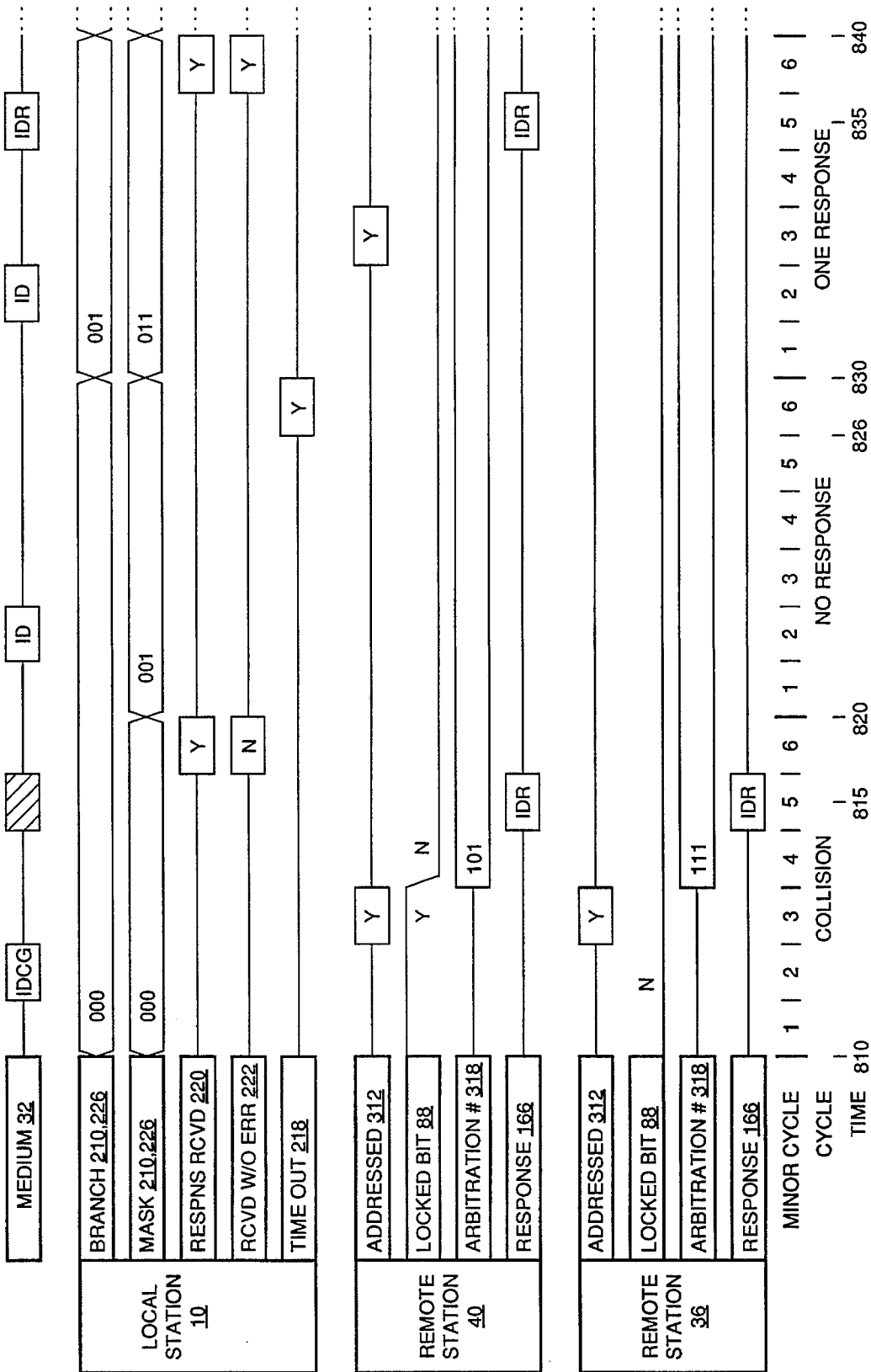


FIG. 13

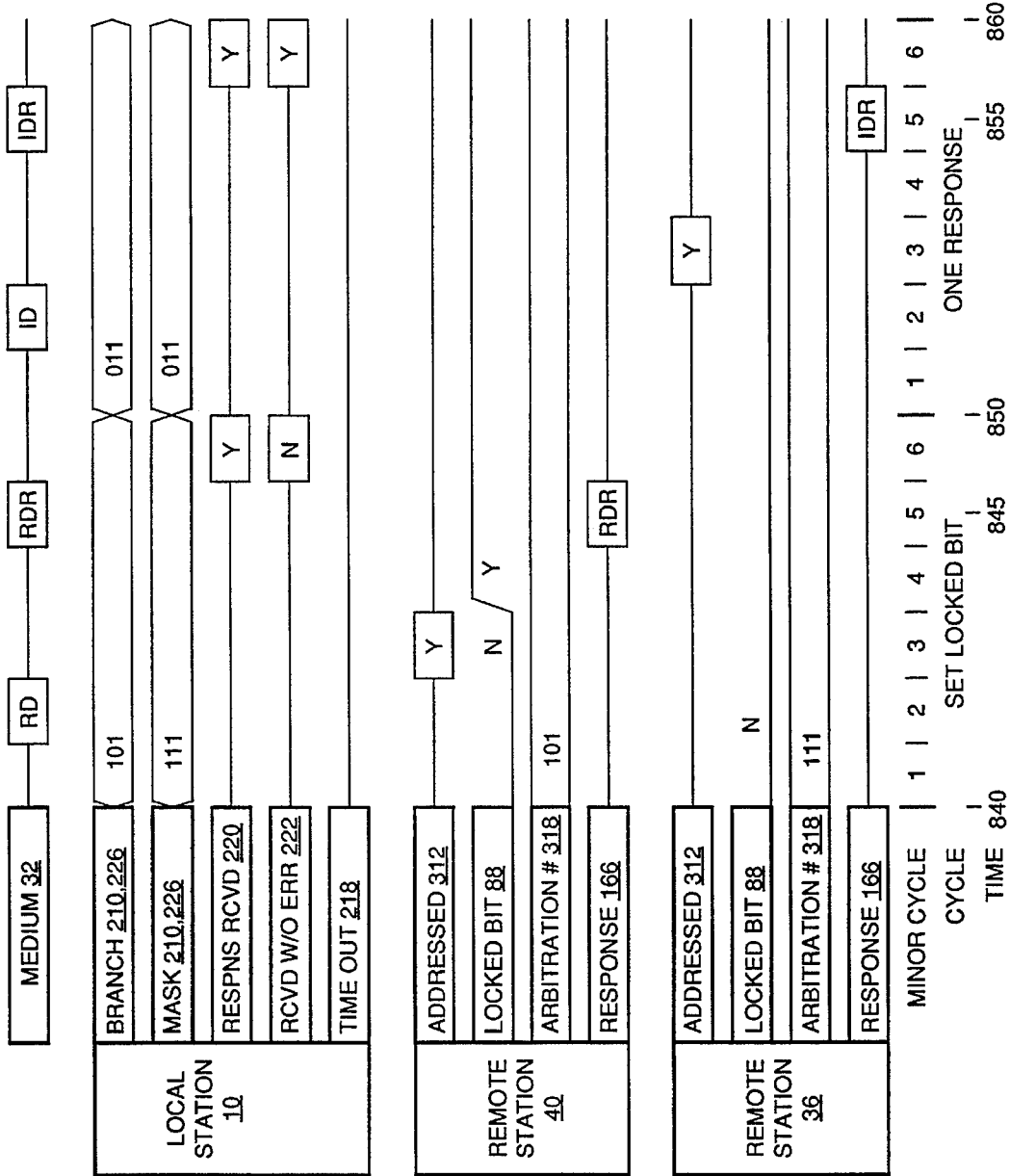


FIG. 14

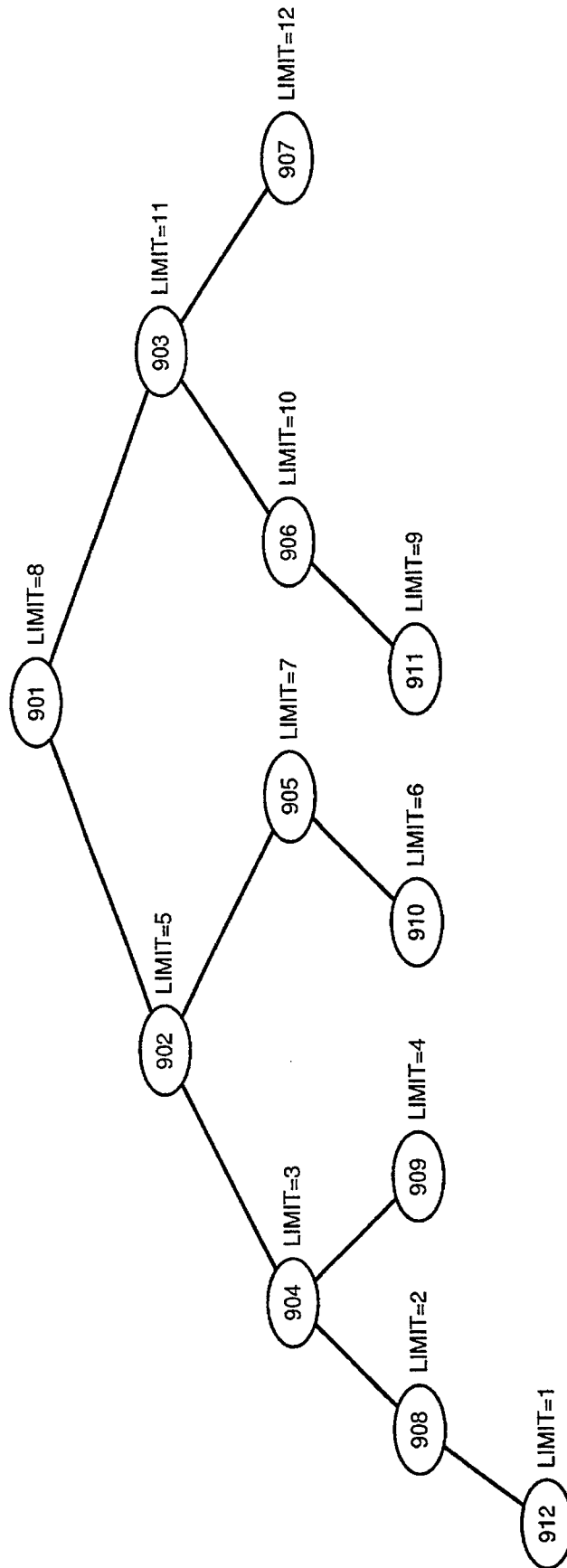


FIG. 15

## DATA COMMUNICATION METHOD USING IDENTIFICATION PROTOCOL

### RELATED APPLICATIONS

The following commonly assigned copending U.S. applications refer to subject matter related to the subject matter of this application: "Radio Frequency Identification Device (RFID) and Method of Manufacture, Including an Electrical Operating System and Method," filed Jul. 7, 1992, Ser. No. 07/899,777; "Electrically Powered Postage Stamp or Mailing or Shipping Label Operative with Radio Frequency (RF) Communication," filed Aug. 11, 1992, Ser. No. 07/928,899; and "Data Communication System Using Identification Protocol," filed December 1992.

### BACKGROUND OF THE INVENTION

This invention relates generally to the field of digital computer data communication. The invention relates more specifically to the data link layer and medium access control of a packet message protocol used in a medium such as radio communication or the like where the quantity and capability of stations requiring use of the medium at a given time cannot be predetermined.

As one example of a communication system, consider a system for handling baggage in an airport terminal. Such a system incorporates radio frequency identification (RFID) between interrogators and responders. In such a system, each baggage tag is a responder station. A station in this sense does not imply stationary location but indicates a device that includes a transceiver for communication. As each bag is transported from its origin to its destination, it sequentially enters areas where interrogation and control functions are to be accomplished by one or more interrogators on all baggage tags in the area within a short period of time. In such an application, as well as in applications where inventory, personnel, animals, packages, samples, mobile stations, and objects must be identified and tracked, there remains a need for communication apparatus and protocol having minimal complexity in circuitry, firmware, and software so that stations can be conveniently equipped and used at practical cost. Whether a communication system is practical depends largely on the system designer's choice of a communication protocol.

A protocol is a method employed uniformly by stations using a common communication medium. Using one method, each station can uniformly determine when it may and when it should not attempt transmission on the medium. By following one protocol, efficient communication can be realized; that is to say that each station's objective for communicating can be accomplished with the least delay and the fewest and shortest messages. Communication efficiency could also be defined as the percentage of time during which only one station is transmitting on the medium. Hence, periods when no station is transmitting and periods when more than one station is transmitting are to be avoided. In the latter case, a collision is said to occur and no intelligible message can be received. The central purpose of a protocol is to provide means for arbitrating between stations that would otherwise cause a collision.

Sophisticated protocols achieve high communication efficiency at the expense of circuit and software complexity required at each station using the protocol. The objectives achieved by communication vary with system design and, for example, may include information transfer, remote sensing and control, and shared computational or shared periph-

eral capability. These objectives are realized by central control of communication or by distributed control. In a system having central access control, permission to use the medium is granted by a central station. When control is distributed, the stations collectively perform a medium access control function to dynamically determine the order in which stations transmit. Of course, interstation communication via the common medium cannot be used to arbitrate. Therefore, the typical protocol includes a set of rules and conditions that each station is constrained in advance to follow. Due to the large number of stations in the modern communication system, random numbers are often employed in arbitration schemes.

Data communication systems have grown in complexity from two-party connections common in the 1970s to systems in the 1990s that interconnect, for example, all employees in a company, all university libraries in a worldwide network, and hundreds of independently managed global networks into a matrix of networks. Such networks are described in "The Matrix, Computer Networks and Conferencing Systems Worldwide" by John S. Quarterman, published by Digital Press, Bedford Mass. (1990), incorporated in full herein by reference. The rapid increase in the number of stations on these networks has been made possible by the adoption of standard protocols. The majority of these standards use distributed access control so that a network can be formed without agreement to a centralized authority. Network standards and access control are further described in "Handbook of Computer-Communications Standards Volume 2 Local Network Standards," by William Stallings, Ph.D., published by Howard W. Sams & Co., Indianapolis, Ind. (1988), incorporated in full herein by reference. After widespread adoption of a standard distributed control protocol, changes to the protocol are costly to implement and, therefore, infrequent. Such change ordinarily would require equipment, firmware, or software changes at each station, and coordination to avoid disruption of important communication. Hence, compatibility with existing networks has led communication systems design away from the development of new protocols, especially new protocols having centralized media access control.

There are many commercial digital communication systems using message packets. Radio communication using message packets are also known. Initial amateur radio packet communication employed a simple acknowledge protocol yielding less than 20% communication efficiency. Improved efficiency was later obtained by adopting a standard distributed medium access control protocol. The description by Quarterman and Stallings of radio networks including the ALOHA network, the Packet Radio Network (PRNET), and the Amateur Radio Packet Radio Network (AMPRNET) is included herein by reference.

Existing popular protocols place a heavy burden of software and circuit complexity upon each station. This burden is illustrated in several standard communication protocols that have been developed for networks of physically connected stations. One ring architecture protocol requires each station to monitor its input channel and retransmit on its output channel the incoming message followed by the station's own message, if any. Another protocol requires each station to transmit only within a time slot beginning at a predetermined time after the station receives a header message broadcast on the medium. A third protocol, commercially known as ethernet, requires each station to perform the following steps:

1. Monitor the medium for inactivity.
2. Generate a random number.

3. Wait a time period based in part on the random number.
4. Attempt to transmit.
5. Detect if a portion of its transmission occurred while another station was transmitting and if so.
6. Begin a subsequent attempt to use the medium by monitoring for inactivity.

Such a method adds to the functional complexity of each station.

In systems having no physical connection between stations, i.e. all stations can receive a broadcast simultaneously, there is a need for a protocol independent of a sequential relationship between stations as in a ring architecture. In systems supporting many thousands of potential stations, time slot assignments are impractical.

Suppose an application requires one station to quickly determine whether all stations currently able to use the medium have been interrogated or commanded in order to accomplish a particular objective through communication. In such an application, ethernet would be unsatisfactory because such a determination depends upon the sum of an indeterminate number of time intervals having randomly selected durations.

Thus, there remains a need for a communication system suited for coordinating the use of a common medium among potentially thousands of stations where no physical connection between stations is desirable and interrogation or control activities must be accomplished in limited time. In addition, there remains a need in some applications to minimize the circuit, firmware, and software complexity required at some stations, perhaps at the expense of complexity at other stations. Without decreasing the complexity, the size and cost per station cannot be reduced to permit new and improved communication systems that employ inexpensive disposable stations such as baggage tags, inventory labels, and the like.

### SUMMARY OF THE INVENTION

A communication system, according to the present invention, employs a common communication medium used by one or more interrogating commander stations and a plurality of responding responder stations. The communication system implements a communication protocol that permits responder stations to be functionally less complex than commander stations. A commander station performs more monitoring and decision making functions than a responder station. Both commander and responder stations generate packets that may contain messages, commands, responses, numbers, coded digital words, and error detection codes. When more than one station transmits on the common medium simultaneously, the resulting interfering transmissions can be detected as an improper transmission at a receiving station, for example, by digital error detection.

According to the present invention, a responder station responds after receiving an error-free message that includes means for the responder station to determine that it has been addressed. When not addressed, the responder station does not respond. A responder station can be addressed individually by including its unique identification number in the message or as a member of a group by including a group designation in the message. Uninterrupted communication between a commander station and a responder station can be conducted in one of two ways. First, when the message transmitted from a commander station includes a group designation to which only one responder station is a member, only one responder station responds. Second, when the

message transmitted from a commander station includes the unique identification of one responder station, only one responder station responds. When the unique identification of responder stations cannot be predetermined, communication, according to one embodiment of the present invention, proceeds as follows.

To establish uninterrupted communication between a commander station and a responder station, the commander station generates and broadcasts a packet containing a command for a first group of responder stations to generate and retain a random number called an arbitration number. When a transmission is intended to be acted upon by a group of receiving stations, the transmission is a type of broadcast called a multicast. Responder stations within the first group are to respond soon after the multicast each giving its arbitration number in a packet containing means for error detection.

After a commander station transmission, several, one, or zero responder stations may simultaneously respond. Following the commander station transmission, when no response is received within a predetermined time, the commander station can ascertain that no responder stations that are members of the first group are currently coupled to the common medium. If an error-free response is received by the commander station, subsequent communication with the responding station can be conducted without interruption when the commander station addresses the responder station as previously discussed. If no message is received within a predetermined time or if a message is received in error (probably caused by several simultaneous responses), the commander station selects a second group and commands the second group to respond.

By broadcasting requests for identification to subsets of the full range of arbitration numbers and checking for a timely error-free response, a commander station can determine the arbitration number of every responder station currently coupled to the common medium.

By requiring initiation of communication by a commander station and substantially immediate response by a responder station, communication systems of the present invention have fewer functional requirements to meet at each responder station. For example, according to the present invention, the responder station is not required to monitor the medium for an opportunity to transmit, monitor the effectiveness of its responses, retransmit responses, or detect retransmitted packets.

According to one aspect, a commander station can establish communication with more than one responder station independent of the operation of other commander stations.

According to another aspect of the present invention, a responder station can sequentially communicate with more than one commander station. Communication systems having multiple commander stations and multiple responder stations are desirable.

According to yet another aspect, a commander station can quickly establish communication with a plurality of responder stations soon after responder stations become coupled to the medium common to that commander station.

According to still another aspect, a responder station can become uncoupled from the medium common to a first group of commander stations and later can be coupled to the medium common to a second group of commander stations without adversely affecting existing or future communication between commander stations and other responder stations.

These and other applications, features, benefits, aspects, and results of the present invention are described below with reference to the accompanying drawings.

## DESCRIPTION OF THE DRAWINGS

FIG. 1 is a functional block diagram of a communication system of the present invention.

FIG. 2 is a functional block diagram of commander station network interface 26 shown in FIG. 1.

FIG. 3 is a functional block diagram of responder station network interface 60 shown in FIG. 1.

FIG. 4 is a diagram of the packet format broadcast by a commander station.

FIG. 5 is a table that describes command opcodes and refers to command formats described in FIG. 6.

FIG. 6 is a diagram of command formats.

FIG. 7 is a diagram of the packet format broadcast by a responder station.

FIG. 8 is a table that describes response mnemonics and refers to response formats described in FIG. 9.

FIG. 9 is a diagram of response formats.

FIG. 10 is a program flow diagram of the protocol followed by a commander station of the present invention.

FIG. 11 is a state diagram of the protocol followed by a responder station of the present invention.

FIG. 12 is a binary tree diagram of BRANCH values and MASK values chosen by a commander station.

FIGS. 13 and 14 are timing diagrams depicting signals in one commander station according to the method of FIG. 10 and of two responder stations according to the method of FIG. 11 as the commander station establishes uninterrupted communication with each responder station.

FIG. 15 is a fibonacci tree diagram of LIMIT values chosen by a commander station.

## DETAILED DESCRIPTION OF THE INVENTION

FIG. 1 is a functional block diagram of communication system 30 of the present invention. In FIG. 1, commander stations 10 and 34 and responder stations 40 and 36 are coupled to common medium 32 by network interfaces 26 and 60, respectively. In practice, a plurality of commander and responder stations would be distributed geographically. The type of medium selected for communication depends on the communication system application; see below for equivalent variations. The embodiment depicted in FIG. 1 illustrates the invention in an application such as airport baggage handling. For this embodiment, the medium is free space through which radio frequency communication are transmissible.

Commander station 10 is designed to achieve a flexible system architecture while incorporating many commercially available components. Commander station 10 includes personal computer system 12 having data and control bus 14 shared by central processor 16, memory 18, and peripheral controllers 20a and 20c. Monitor 22, disk drive 24, and network interface 26 connect to individual peripheral controllers 20a-20c via connecting signals 28a-28c, respectively. Network interface 26 is coupled to common medium 32. Network interface 26 could be implemented in a chassis separate from the chassis of personal computer system 12 or equivalently could be implemented in combination with the functions of the network interface peripheral controller 20c for connection directly to data and control bus 14.

The configuration of commander station 10 illustrates several advantages. Communication system configuration and operation are largely dictated by software loaded via

disk drive 24, stored in memory 18, and performed by central processor 16. Disk drive 24, memory 18, central processor 16, as well as monitor 22 and peripheral controllers 20a-20c are all conventional, general purpose, and readily available apparatus. Therefore, additional functions and changes to communication system 30 can be made in software with little or no mechanical changes to commander station 10. The operation of commander station software will be discussed below.

Responder station 40 is designed for minimal circuitry to achieve, among other things, small size and low power consumption. Small size permits convenient use, for example, as a baggage tag. Low power consumption permits further size reduction and reduces manufacturing and operating costs. Small size and low manufacturing costs combine to permit implementing responder station 40 as a convenient, dispensable, throw-away item such as a baggage tag, package label, or the like.

In essence, microsequencer 42 forms the core of responder station 40. Microsequencer 42 is a read only memory that produces data signal 48 in response to address signal 44. In operation, a value is presented as address signal 44 once every period of clock signal 46. Data signal 48 from microsequencer 42 is stored in state register 50 once every period of clock signal 46. The output of state register 50 is state signal 52, which forms control bus 54. Control bus 54 causes register transfer operations to be described below. A portion of state signal 52 defines a portion of address signal 44. Thus, a sequence of state transitions occurs in synchronism with clock signal 46 as defined by the internal operation of microsequencer 42 and other signals together comprising address signal 44, as will be described below. A state transition diagram is also discussed below. In the typical microsequencer, internal multiplexers reduce the range of read only memory addresses that would otherwise be required. Microsequencer 42 is of a class of devices described by Charles Belove in "Handbook of Modern Electronics and Electrical Engineering," pp 2135-2142, published by John Wiley & Sons, New York, NY (1986), incorporated herein by reference; and by Y. Chu in "Computer Organization and Microprogramming," published by Prentice-Hall, Englewood Cliffs, N.J. (1972) incorporated in full herein by reference.

Network interface 60 of responder station 40 is coupled to common medium 32 in a way similar to the coupling of network interface 26 of commander station 10 to common medium 32. Network interface 60 connects to state register 50 to supply clock signal 46. Network interface 60 connects to control bus 54 so that send and receive operations are directed in part by state signal 52. When microsequencer 42 is in an appropriate state, data received by network interface 60 is transferred from network interface 60 to command register 56 by data bus 62 in conjunction with load signal 58. Command register output 59 defines a portion of address signal 44. Network interface 60 also connects via data bus 62 to memory 64, register array 66, flag register 84, and random number generator 90 for transfer of data between these function blocks, for storage of received data, and for recall of data to be sent.

In one embodiment, data bus 62 is byte-wide. Network interface 60 converts received data from serial to byte-parallel organization. The several devices that connect to data bus 62 make a byte-parallel connection.

In another embodiment, data bus 62 is bit-serial. Control bus 54, in such an embodiment, includes serial clock signal (not shown). Register transfer among network interface 60,

register array **66**, memory **64**, flag register **84**, and random number generator **90** are accomplished in bit-serial fashion with appropriate electrical interfaces known in the art.

In yet another embodiment, a combination of serial and parallel data paths are implemented. The system designer's choices of serial or parallel as well as the number of bits per register transfer operation, depends on factors including system and device timing limitations, noise immunity, power dissipation, device size, topology, and layout constraints.

Memory **64** connects to read/write signal **68** and memory address signal **70** which are part of control bus **54**. Memory **64** is used to store values for responder station identification and data related to the communication system application. For example, when a responder station is used as an airline baggage tag, postal mailing label, or inventory control tag, memory **64** would store data describing a destination for the item to which the tag is attached.

Register array **66** performs functions similar to a multi-port memory. Register array **66** connects to arithmetic-logic unit (ALU) **72** for the presentation of operand signals **74** and **76**, and storage of result signal **78**. Operand and result signals are multi-bit digital signals for arithmetic operations such as addition, bit-wise parallel logical operations such as logical-AND, and bit-wise serial operations such as shift-left. Register array **66** connects to control bus **54** so that registers to be coupled to operand and result signals are selected and stored according to state signal **52**.

In addition to the connections already described, a portion of control bus **54** connects to ALU **72** to supply opcode signal **80** to ALU **72**. Opcode signal **80** selects one of a plurality of possible operations to be conducted by ALU **72**. When an equality comparison has been selected by opcode signal **80** and operand signals **74** and **76** are bit-wise identical, A=B signal **82** is asserted by ALU **72**. A=B signal **82** defines a portion of address signal **44**.

Control bus **54** connects to individual bits arranged in flag register **84**. Addressed-bit **86**, part of flag register **84**, is set under control of state signal **52** to indicate whether responder station **40** has been addressed in a received command message. Locked-bit **88**, also part of flag register **84**, is set under control of state signal **52** to indicate whether responder station **40** should ignore messages from a commander station because responder station **40** has already announced its identification to a commander station. The significance of addressed-bit **86** and locked-bit **88** will become more readily apparent in the description below.

Random number generator **90** connects to control bus **54** and data bus **62** for transferring a random number of a predetermined precision to register array **66**. When retained in register array **66**, the random number is called an ARBITRATION NUMBER whose function will be discussed below. Circuit techniques for generating a random number in digital format are well known and described, for example, by H. F. Murray in "General Approach for Generating Natural Random Variables", IEEE Transactions on Computers, Vol. C-19, No. 12, pp 1210-1213, December 1970, incorporated herein by reference. In one embodiment, random number generator **90** is similar to an integrated circuit implementation described by Alan Folmsbee, et. al., in their article, "128K EPROM with Encrypted Read Access", published in the Digest of Technical Papers IEEE International Solid-State Circuits Conference pp 46-47 and 103, by Lewis Winner, Coral Gables, Fla., 1985, incorporated herein by reference.

FIG. 2 is a functional block diagram of network interface **26** shown in FIG. 1. Within network interface **26**, connecting

signal **28c** couple to output buffer **110**. Byte-parallel loading of output buffer **110** is accomplished by network interface peripheral controller **20c** shown in FIG. 1. Bytes are removed from output buffer **110** by transmitter logic **112** to accomplish several processing objectives. In one embodiment, a 5-bit cyclic redundancy check (CRC) code is joined to each 8-bit byte to form a 13-bit word. Redundancy, provided by the 5-bit code, facilitates error detection and limited error correction by responder station **40**. Table 1 describes the format of the 13-bit word and includes a description of the CRC code used in one embodiment. Suitable CRC encoder and decoder circuits used in transmitter logic **112** and receiver logic **178** are described in detail in "Error Control Coding: Fundamentals and Applications," by Shu Lin and Daniel J. Costello Jr., Prentice-Hall Englewood Cliffs, N.J. 1983, pp 62-94. Transmitter logic **112** also generates transmit serial bit stream **114** which includes a message preamble bit stream, one or more successive 13-bit words from output buffer **110**, and a postamble bit stream. When permitted by OK-to-transmit signal **116**, transmit serial bit stream **114** is presented to transmitter **118**. Transmitter **118** in one embodiment produces a radio frequency transmit signal **120** by modulation and couples that signal to antenna **122**. Appropriate modulation methods depend on the communication medium.

TABLE 1

Bit Order of Transmission	D7,D6, . . . D0,P4,P3, . . . P0 (P0 transmitted first)
CRC Generation Equations	P0=D1+D2+D5+D7 P1=D1+D3+D4+D6 P2=D0+D2+D3+D6+D7 P3=D0+D4+D5+D6+D7 P4=D0+D1+D2+D3+D4+D5

For a communication system for airport baggage handling, modulation includes, for example, spread spectrum modulation having pseudo noise characteristics. Other techniques for transmitter design appropriate to radio transmission and other media will be readily apparent to those skilled in the arts applicable to communication on a particular medium. A power level of approximately 1 watt is sufficient to excite responder station network interface **60** at distances and noise levels typically required for a communication system for airport baggage handling.

Receiver **124** is coupled to antenna **122** for amplifying and filtering radio frequency received signal **126**. Receiver **124** derives OK-to-transmit signal **116** from power level measurements on received signal **126** and provides signal **116** to transmitter logic **112**. Although responder station network interface **60** need not generate a transmitted signal using the same modulation technique employed in transmitter logic **112** and transmitter **118**, a common method is preferred, for example, in order to minimize circuitry in responder station network interface **60**. Thus, receiver **124** removes the carrier signal and other artifacts of modulation generated by responder station network interface **60** in one embodiment by synchronizing with the spread spectrum signal and removing pseudo noise characteristics through known detection and filtering methods. Resulting received serial bit stream **128** is coupled to receiver logic **130**, in one embodiment, for determining whether a proper message has been received and for decomposing the message into successive 8-bit bytes. The method and circuitry required to determine whether a proper message has been received depend on the redundancy that responder station network interface **60** incorporates into received serial bit stream **128**.

For a communication system for airport baggage handling, responder station **40** may transmit at a power level of



1 milliwatt or less. Multiple and more sophisticated error detection schemes transmitted from responder station 40 will extend the limit of physical separation between commander station 10 and responder station 40. Error detection schemes are well known. Such schemes may also permit reliable communication in environments having substantial noise levels. On the other hand, limits to the complexity and power consumption of responder station 40 may limit the extent of encoding circuitry therein. Where responder station network interface 60 facilitates one or more particular error detection schemes, receiver logic 128 decomposes, decodes, detects, and to a limited extent corrects errors in received serial bit stream 128. In one embodiment, receiver logic 130 determines proper-message-received signal 132 by decoding Viterbi encoding using model Q1601 decoder available from Qualcomm, Inc., San Diego, Calif. according to Qualcomm application notes and the parameters: Rate  $R=1/2$ , Generating Functions  $G0=171$ (octal) and  $G1=133$ (octal), and Constraint Length  $K=7$ . Receiver logic 130 also performs serial to parallel conversion to produce successive 8-bit bytes which are stored in input buffer 134.

Network interface peripheral controller 20c, responsive to OK-to-transmit signal 116 and proper-message-received signal 132, generates signals on data and control bus 14 from which central processor 16 can determine whether a proper message has been received as of a predetermined time after transmission and if not, whether no message was received. Other control signals, known in the art and not shown, are generated and sensed to orchestrate the loading of output buffer 110 and the unloading of input buffer 134 under control of central processor 16.

FIG. 3 is a functional block diagram of responder station network interface 60 shown in FIG. 1. The configuration illustrated in FIG. 3 performs functions similar to those already described above for commander station network interface 26. Differences between the two serve primarily to limit the complexity of responder station circuitry. Data bus 62 connects to transmitter logic 160. When directed by microsequencer 42 via signals on control bus 54, transmitter logic 160 generates a message preamble bit stream. Then, for each bit of each word read from memory 64 or register array 66, transmitter logic 160 develops a Viterbi code. Functional descriptions suitable for designing circuits or computer programs for generating Viterbi and similar convolutional codes are explained in "Error Control Coding: Fundamentals and Applications," by Shu Lin and Daniel J. Costello Jr., Prentice-Hall Englewood Cliffs, N.J. 1983, pp 287-456; and "Error-Correction Coding for Digital Communications", by George C. Clark, Jr. and J. Bibb Cain, Plenum Press, New York, N.Y. 1981, pp 227-266. Message signal 162 presents the codes to transmitter 164. Following the last code, transmitter logic 160 generates a message postamble bit stream. Transmitter 164 modulates message signal 162 in a way compatible with receiver 124 and receiver logic 130. The resulting transmit radio frequency signal 166 is coupled to antenna 168. Redundancy, provided by the Viterbi codes, facilitates error detection and limited error correction when the message is received at commander station 10.

Receiver 170 is coupled to antenna 168 for amplifying and filtering radio frequency received signal 172. Receiver 170 derives wake-up signal 174 from power-level measurements on received signal 172 and provides wake-up signal 174 to power control and restart circuits not shown. In a communication system for airport baggage handling, non-critical circuitry in responder station 60 is powered by battery only after the preamble of a packet has been

detected. Receiver 170 also removes the carrier signal and other artifacts of modulation generated by commander station network interface 26 in one embodiment by synchronizing with the spread spectrum signal and removing pseudo noise characteristics through detection, demodulation, decoding, and filtering methods known in the radio communication arts. Resulting received serial bit stream 176 is coupled to receiver logic 178. In one embodiment, receiver logic 178 performs several functions: Determining whether a proper byte has been received, consequently generating improper-byte-received signal 180, and decomposing the packet into successive 8-bit bytes forming received message signal 182. The method and circuitry required to determine whether a proper byte has been received depend on the redundancy that commander station network interface 26 incorporates into transmit serial bit stream 114. Receiver logic 178 detects the first byte of a command and in response generates load signal 58. Clock signal 46 is also generated by receiver logic 178 to drive state register 50.

Microsequencer 42 and network interface 60 cooperate via control bus 54 which includes improper-byte-received signal 180. Other signals included in control bus 54, known in the art and not shown, orchestrate transfer of bytes between memory 64, register array 66, transmitter logic 160, and receiver logic 178. If an improper byte is received, as indicated by improper-byte-received signal 180, microsequencer 42 responds by reverting to an idle state and ignoring incoming bytes until another command is received.

For a detailed description of suitable circuits of the type that can be used for transmitters 118 and 164 and receivers 124 and 170, implemented in spread spectrum technologies, see U.S. Pat. No. 5,121,407 by Partyka et al., incorporated herein by reference.

FIG. 4 is a diagram of the packet format sent by commander station 10 to responder station 40. Each command packet 140 includes, in order of transmission, a preamble followed by a command followed by a postamble. The preamble and postamble are designed for synchronizing a transmitter circuit and a receiver circuit for a particular packet. In one embodiment, the preamble bit stream comprises 768 '1' bits followed by a 7-bit Barker code of '0001101'. In one embodiment, the postamble comprises a 7-bit Barker code of '1110010'.

In one embodiment, each bit of the command format is modulated using a pseudo noise (PN) sequence for direct sequence spread spectrum communication. The sequence is generated in part by a linear feedback shift register (LFSR) of the form [5,2]. In this form, the input to the first of five registers is the result of combining the output of register 5 by exclusive-OR with the output of register 2. The generator in this embodiment has 32 states so that the 1 and 0 states occur with equal probability. Since the LFSR generates only 31 states, an additional state is inserted by support circuitry. For a detailed description of a suitable PN modulator circuit of the type employed in transmitter 118 see "Spread Spectrum Systems", by R. C. Dixon, published by John Wiley and Sons, Inc. 1984 pp 15-28 and 56-151 incorporated herein by reference. Suitable demodulator techniques and circuits (of the type used in receiver 170 to recover the response format) are also described at pages 153-290 incorporated herein by reference.

FIG. 5 is a table that describes several commands and refers to command formats described in FIG. 6. As shown in FIG. 6, each command begins with an opcode and has one of four formats varying in length from 3 bytes to 258 bytes. Opcode values were selected to facilitate accurate decoding

and obtain high noise immunity. Each byte is an 8-bit word as it would appear in output buffer 110 and on received message signal 182. The opcode hexadecimal value is stored on receipt in command register 56. Bytes following the opcode have the following meanings. MASK and BRANCH as used in format 142 are binary numbers chosen by a commander station to specify a group of responder stations that should act on the command and should reply. LOCAL ID in format 142 is a unique identification number assigned, for example, by the communication system installer to each commander station 10, 34 coupled to common medium 32. Responder stations 36, 40 coupled to common medium 32 can then direct a response to one of several commander stations 10, 34 by, for example, including a particular LOCAL ID in each response. When one commander station chooses to specify only one responder station that should act on a command and should reply, that commander station includes in its command an ARBITRATION NUMBER as in formats 144 and 146 identifying the responder station. An ARBITRATION NUMBER is a short value, for example 1 byte, chosen for self identification by a responder station. On the other hand, A TAG, as in format 146, is a long value, for example 8 bytes, assigned by a communication system designer at the time a responder station is manufactured or commissioned. The ARBITRATION NUMBER distinguishes responder stations when coupled simultaneously with at least one commander station to a common medium. However, the TAG, distinguishes responder stations throughout the life of the communication system application. Finally, DATA in format 146 includes some or all of the contents for any or all devices including memory 64, register array 66, flag register 84, and random number generator 90.

FIG. 7 is a diagram of the packet format sent by responder station 60 to commander station 10. Each response packet 190 includes, in order of transmission, a preamble followed by a response followed by a postamble. The preamble and postamble are designed for synchronizing a transmitter circuit and a receiver circuit for a particular packet. In one embodiment, the preamble bit stream comprises 768 '1' bits followed by a 7-bit Barker code of '0001101'. In one embodiment, the postamble comprises a 7-bit Barker code of '1110010'.

In one embodiment, each bit of the response format is modulated using a pseudo noise (PN) sequence for direct sequence spread spectrum communications. The sequence is generated in part by a linear feedback shift register (LFSR) of the form [6,1] or [8,4,3,2] for either a 64 chip sequence or a 256 chip sequence respectively. In the form [6,1], the input to the first of six registers is the result of combining the output of register 6 by exclusive-OR with the output of register 1. Similarly, for the [8,4,3,2] form, the input to the first of eight registers is the result of the exclusive-OR of the outputs of registers 8, 4, 3, and 2. The 64 chip sequence requires less time for signal synchronization than the 256 chip sequence; however, the latter provides better performance in systems having poor signal to noise ratio. The generator in this embodiment has an even binary multiple of states, so that the 1 and 0 states occur with equal probability. Since the LFSR generates one less state, an additional state is inserted by support circuitry. For a detailed description of a suitable PN modulator circuit of the type employed in transmitter 164, see "Spread Spectrum Systems", by R. C. Dixon, published by John Wiley and Sons, Inc. 1984 pp 15-28 and 56-151 incorporated herein by reference. Suitable demodulator techniques and circuits of the type used in receiver 124 to recover the response format are also described at pages 153-290 incorporated herein by reference.

FIG. 8 is a table that describes several responses and refers to response formats described in FIG. 9. As shown in FIG. 9, response formats 192-196 include LOCAL ID, ARBITRATION NUMBER, and TAG, which have the meanings already described above. By including LOCAL ID and ARBITRATION NUMBER in each response, in cooperation with locked bit 88 one responder station can respond unambiguously to one commander station in the presence of a plurality of commander and responder stations. The INVERTED ARBITRATION NUMBER in format 192 is the binary ones-complement of the ARBITRATION NUMBER and is included for increased accuracy of communication. REVISION in format 192 is a one-byte value set by a communication system developer at the time of manufacture or commissioning of a responder station. REVISION represents the responder station configuration and connotes its capability. STATUS in format 196 is a one-byte code chosen by responder station 40 to convey current conditions of important system events such as low battery, uncorrectable data received, write protection, and similar information which may indicate to commander station 10 that communication should be repeated or abandoned. DATA in response format 194 includes some or all of the contents of any or all devices including memory 64, register array 66, flag register 84, or random number generator 90.

A communication system, according to the present invention, includes commander and responder stations that adhere to a method of communicating called a protocol. In general, the protocol of the present invention places different requirements on a commander station than on a responder station. Thus, there is a commander station method (FIG. 10) and a responder station method (FIG. 11). These methods together implement the communication system protocol.

Operation according to the present invention produces the following characteristic effects at the system level. First, a commander station will not begin transmitting during the transmission by another commander station or by a responder station. Operation, according to the present invention, does not prevent more than one commander station from beginning transmission simultaneously; however, it is feasible to couple commander stations to a second medium or to constrain commander stations to a second or expanded protocol on common medium 32. For example, commander stations 10 and 34 include personal computer system 12, which can be augmented with a peripheral controller for operation over ethernet. Communication over the second medium can be used to prevent simultaneous broadcast over common medium 32. For example, a second protocol on common medium 32 may include operator action to assign time slots, back off delays, or similar means for media access whether central or distributed. Several embodiments for these means for media access have been described by Stallings in his work already incorporated by reference above.

Second, a responder station will not transmit unless it has first received a command to which it determines it must respond. The response is made within a predetermined time immediately following receipt of the command.

Third, a commander station can form a command in a manner calling for all, more than one, or one responder station to respond. An important object of the communication system protocol in a communication system of the present invention, i.e. uninterrupted communication, is achieved after a commander station determines how to cause only one responder station to respond. The program flow diagram of FIG. 10 and the state diagram of FIG. 11 describe how uninterrupted communication between one commander

station and each responder station is achieved when a plurality of commander stations and a plurality of responder stations are simultaneously coupled to a common medium.

FIG. 10 is a program flow diagram of the protocol followed by a commander station of the present invention. A practical example of a communication system will be used to describe the flow diagram.

In a communication system for airport baggage handling the quantity and identity of responder stations within the radio communication range of a commander station varies over time. A commander station may be at a fixed operator station within radio range of a moving belt conveying baggage toward a Y-junction of conveyors. When baggage tags are constructed as responder stations and when each tag has destination information stored in memory 64, the commander station, through communication with each baggage tag, can control the routing of each bag through the junction onto one of two conveyors. Assume that each responder station also has information in memory 64 describing its sequential position on the conveyor. Such a sequence number could be a date and time of day when the bag passed through a chute upstream of the commander station.

As a group of bags approaches the commander station, the commander station has a fixed amount of time to determine the identity of each responder station, in order to establish uninterrupted communication. For proper baggage handling, the commander station must routinely and repeatedly identify all bags on the conveyor. To do so, at FIG. 10 block 210, commander station 10 specifies a group of responder station addresses by choosing values for BRANCH and MASK. BRANCH and MASK values are determined in a manner to be explained by reference to FIG. 12 below. In one embodiment, the initial group specification, i.e. BRANCH and MASK values, would specify all possible responder stations. Commander station 10 at block 212 generates an "identify, clear, and generate" (IDCG) command having a format according to FIGS. 4, 5, and 6. When the media is clear to broadcast, block 214, as indicated by OK-to-transmit signal 116, the IDCG message is broadcast, block 216. An IDCG message causes each responder station that is a member of the group to clear locked-bit 88, generate a random number and retain it as its ARBITRATION NUMBER, and broadcast a response. The responder station's reactions to ID, IDG, IDC, and IDCG commands are explained further in reference to FIG. 11 below.

Commander station 10 now loops through blocks 220 and 218 for a response to be received as indicated by OK-to-transmit signal 116 or a time out elapsed condition. If a response was received, as indicated by a false state of OK-to-transmit signal 116, commander station 10 at block 222 determines whether a collision occurred, as indicated by a false state of proper-message-received signal 132. If commander station 10 determines that a collision occurred, it will determine at block 224 whether all possible members of the initial group of responder station addresses specified at block 210 have been addressed in an ID, IDG, IDC, or IDCG command. How this determination is made will be further explained with reference to FIG. 12 below. If all subgroups have not been tried, the commander station again specifies a group of responder station addresses, for example, a subgroup or disjoint group of a prior group. At block 228 commander station 10 generates an ID command according to FIGS. 4, 5, and 6 and continues the method from block 214.

If, at block 218, a predetermined time elapsed without a false condition appearing on OK-to-transmit signal 116,

commander station 10 concludes that no response was transmitted and continues the method at block 224.

If, at block 222, the proper-message-received signal is true, then commander station 10 concludes that only one responder station responded. At block 230, commander station 10 determines and validates the responding responder station's ARBITRATION NUMBER according to response format 192 using ARBITRATION NUMBER and INVERTED ARBITRATION NUMBER. According to a particular system communication objective, commander station 10 now selects a command from FIG. 5 which will cause the responder station to set its locked-bit 88. For determining baggage destination and positional sequence on the conveyor, commander station 10 could select RD. Using the appropriate command format shown in FIG. 6, commander station 10 generates a message at block 232, loops until the OK-to-transmit signal indicates that the medium is clear to broadcast at block 234, then broadcasts the command at block 236. Commander station 10 again awaits a proper response message by looping at block 238 through block 240. If a predetermined time elapses at block 240, commander station 10 continues the method at block 234. If a response is received without error at block 244, as indicated by proper-message-received signal 132, then two party uninterrupted communication between commander station 10 and one responder station 60 has been established. Further communication may be required, as indicated by the STATUS code in the received response format 192 or to accomplish other system communication objectives.

It is possible at block 224 for the commander station to determine that no further subgroup, super group, or disjoint group remains to be commanded using the ID command. Suppose, for example, that all practical values of BRANCH and MASK have been used. If the immediately preceding broadcast at block 216 elicited no response at block 218, then commander station 10 can conclude that all responder stations have been identified. Otherwise, at block 248, commander station 10 generates an identify and generate command (IDG) according to the format in FIGS. 4, 5, and 6 using the same group that was specified in block 210. Commander station 10 continues the method at block 214.

Although the same group is specified, a responder station that has been identified at block 244 will not respond, since its locked-bit 88 has been set. Collisions are less likely to occur with each pass through the loop from block 214 to block 248 because a smaller number of responder stations can respond. Hence, the method of FIG. 10 converges toward identifying all responder stations. The communication system designer must select the precision of BRANCH and MASK values to assure conversion within system time allowances, for example, 8-bit BRANCH and MASK values are compatible with conveyor speeds and radio ranges needed for airport baggage handling.

FIG. 11 is a state diagram of the protocol followed by a responder station of the present invention. Responder station 40, begins in idle state 310 when power is applied or restored according to wake-up signal 174. In part, the idle state is indicated by contents of command register 56 not corresponding to a valid command. The idle state is re-entered to interrupt command processing when improper-byte-received signal 180 is raised by receiver logic 178. A valid command loaded into command register 56 causes state transition to address check state 312.

In address check state 312, microsequencer 42 determines whether responder station 40 has been addressed by one of two methods. First, if the command conforms to format 142,

the responder station is addressed when the result of ARBITRATION NUMBER logically ANDed with MASK is bitwise identical to BRANCH. ARBITRATION NUMBER is the current contents of a particular register in register array 66. MASK and BRANCH are values received in the command and stored in register array 66. Logical operations and comparisons are performed by ALU 72 which produces A=B signal 82. If A=B signal 82 is not asserted, state 314 is entered. Responder station 40 may remain in state 314 until a predetermined time elapses. Responder station 10 re-enters idle state 310, after the predetermined time elapses.

To illustrate the importance of such a delay, suppose that commander and responder stations employed radio transceivers for network interfaces 60 and 26. Then, suppose responder station 40 is within range of two commander stations 10 and 34, but commander stations 10 and 34 are out of range of each other. When commander stations 10 and 34 validly produce back to back commands, the delay interposed by state 314 prevents responder station 40 (not addressed by commander station 10 in the first occurring command) from responding to commander station 34 in the second occurring command. Without the delay, a collision could occur that may confuse commander station 10.

A second way to determine whether responder station 40 has been addressed applies for commands having formats 194 and 196. Accordingly, responder station 40 is addressed when ARBITRATION NUMBER, retained in register array 66, is bit-wise identical to ARBITRATION NUMBER as received in the command. Comparison is performed by ALU 72 which produces A=B signal 82. If A=B signal 82 is not asserted, state 314 is entered as already described. Otherwise, transition is made to decode state 316.

Decode state 316 follows address check state 312 in response to A=B signal 82. If the command opcode is not recognized then no response state 314 is entered. For some commands, a further condition such as the state of locked-bit 88, if not met, will cause the command to be treated as not recognized. Each recognized command opcode causes microcode execution to begin at a section of microcode for the purpose of directing microsequencer operations to process the particular received command. Four commands are illustrated as separate states 318 through 324 and other commands are illustrated in summary by pseudo state 326.

When the opcode for command IDCG has been received, state is entered for identify, clear, and generate operations. An IDR response (according to FIGS. 8 and 9) is selected, locked-bit 88 is cleared, the content of random number generator 90 is stored in register array 66 as ARBITRATION NUMBER, and transition is made to state 328.

When the opcode for command IDG has been received and locked-bit 88 is not set, state 320 is entered for identify, and generate operations. An IDR response is selected and a new ARBITRATION NUMBER is generated as already described for state 318. Transition is then made to state 328.

When the opcode for command IDC has been received, state is entered for identify and clear operations. An IDR response is selected and locked-bit 88 is cleared. Transition is then made to state 328.

When the opcode for command ID is received and locked-bit is not set, state 324 is entered for an identify operation. An IDR response is selected. Transition is then made to state 328.

When the opcode for other commands (including RD and WD) is received, locked-bit 88 may be set and other functions may be performed. Other functions may include writing data to memory 64, writing data to register array 66,

altering the state of registers including flag register 84, and other operations controlling responder station configuration and operation. Transition is then made to state 328.

Upon transition to state 328, the response selected by a prior state is generated according to FIGS. 7, 8, and 9 and broadcast. In one embodiment, the response is broadcast as it is being generated. Transition to idle state 310 is made, after broadcasting the response. Note that responder station 40 does not wait for clear medium prior to broadcasting the response. According to the present invention, collision detection by responder stations is not necessary to accomplish uninterrupted communication.

The ARBITRATION NUMBER generated by responder station 40 and the BRANCH and MASK numbers chosen by commander station 10 operate to establish uninterrupted communication. We now turn to a further explanation of the method used by commander station 10 to choose BRANCH and MASK values.

FIG. 12 is a binary tree diagram of BRANCH values and MASK values chosen by a commander station. A tree is a type of graphic representation. There are several types of trees known in mathematics and computer science. The tree depicted is a binary tree where a node can have two branches, shown descending left and right from a node. Each node corresponds to a unique combination of values for BRANCH and MASK, which are n-bit binary numbers having the same precision. As illustrated, BRANCH and MASK are 3-bit binary numbers. In a communication system for airport baggage handling, 8-bit numbers would be used. The precision employed for BRANCH and MASK must be identical to the precision selected for ARBITRATION NUMBER generated by responder station 40.

Recall that responder station 40 uses the expression ARBITRATION NUMBER AND MASK=BRANCH to determine if it is addressed, where ARBITRATION NUMBER is the value retained in register array 66 from random number generator 90. When MASK is 0 and BRANCH is 0 all values of ARBITRATION NUMBER satisfy the expression, i.e. all responder stations coupled to common medium 32 conclude they are addressed. On the other hand, if MASK has a '1' bit in every position, then the expression is satisfied for only one value of ARBITRATION NUMBER.

When MASK is arranged with '0' and '1' bits, the expression is satisfied by a group of values for ARBITRATION NUMBER, and potentially a group of responder stations could conclude they are addressed. Note for a responder station to be addressed, BRANCH at bit position 'q' must be '0' when MASK at bit position 'q' is '0', for all values of 'q'. When MASK at bit position 'q' is '1', BRANCH can take two values for that bit position which correspond to the left and right branches of a binary tree.

At the first level of the tree, nodes 702 and 703, MASK is '1' in bit position 'r'. The corresponding bit position of BRANCH is '0' at node 702 and '1' at node 703. At the second level of the tree, nodes 704 through 707, MASK is '1' at bit positions 'r' and 's'. For example, the value for BRANCH at node 707 is the parent node BRANCH value (001 at node 703) modified by forcing a '1' (for the right-hand branch) at bit position 's', hence 011. In like manner, the BRANCH and MASK values for any node in the tree can be determined. In FIG. 12 MASK bit positions have been given in an order right to left. Any other order of bit positions would be equivalent. Methods for choosing first and subsequent values for BRANCH and MASK can now be explained in terms of traversing from node to node on a binary tree.

When commander station 10 broadcasts a request for identification (an ID, IDC, IDG, or IDCG command) one of three events can occur. BRANCH and MASK values given at a particular node that represents a first group of responder stations. First, commander station 10 could receive no response from which it could conclude that no responder station in the first group is currently coupled to the common medium 32. Second, a proper response could be received. From that event, commander station 10 could conclude that only one responder station in the first group is currently coupled to common medium 32. Third, from an improper response received, commander station 10 could conclude that a collision of more than one response occurred. An improper response could be caused by, for example, weak coupling, high noise levels, or weak received signals. However, these causes can be treated in the same way as a collision to simplify the commander station protocol without substantially degrading system performance for applications including airport baggage handling. Therefore, an improper response simply merits further search.

An efficient search for the identity of each of several responding responder stations is equivalent to an efficient search for the leaves of a binary tree. A leaf is a node having no further branches. When use of the values for BRANCH and MASK at a node produces no collision, the node is a leaf. Tree search methods are easily implemented using known computer programming methods.

Tree search methods are essentially of two types, breadth first and depth first. A particular communication system application may use one method or the other to optimize commander station computing time and memory space objectives. An explanation of these methods using the programming language PASCAL is given by E. Horowitz and S. Sahni in "Fundamentals of Data Structures in PASCAL" pp 203-265 and 326-332 published by Computer Science Press Inc., Rockville, Md. (1984), incorporated herein by reference.

Suppose that two responder stations 40 and 36 and one commander station 10 are coupled to common medium 32. The binary tree in FIG. 12 illustrates a sequence of BRANCH and MASK values used by commander station 10 to identify responder stations. Timing diagrams in FIGS. 13 and 14 illustrate the same sequence showing decisions at commander station 10 decision blocks (according to the commander station method of FIG. 10) and responder station control signals (according to the responder station method of FIG. 11) as commander station 10 establishes uninterrupted communication with each responder station.

Beginning at FIG. 10 block 210, FIG. 12 node 701, and FIG. 13 time 810, commander station 10 chooses BRANCH=000 and MASK=000, calling for all responder stations to respond. At time 815, responder station 40 has determined that it is addressed, has cleared its locked-bit 88, has generated ARBITRATION NUMBER 101, and has begun transmitting response IDR. Simultaneously, responder station 36 has determined that it has been addressed, has generated ARBITRATION NUMBER 111, and has begun transmitting response IDR. Also, at time 815, simultaneous transmissions collide on common medium 32.

At time 820, commander station 10 at block 226 chooses node 702 having BRANCH=000 and MASK=001. Responder station 40 is not addressed because ARBITRATION NUMBER (101) ANDed with MASK (001) yields 001 which is not equal to BRANCH (000). Similarly, responder station 36 is not addressed. Neither station responds. At time 826, time out at block 218 occurs.

At time 830 and block 226, a third group of responder station addresses is chosen. From FIG. 12 the appropriate group is specified by traversing the tree according to a search method. If a breadth first search were used, all nodes at the same level would be visited before testing at a deeper level. Hence, node 703 would be next. If a depth first search were used, search would proceed upward from node 702 (because it is a leaf) and then downward from the first node having an untested branch. Hence, up to node 701 and down to node 703. As a refinement to either method, node 703 can be skipped because a collision at node 701 and no response at node 702 implies a collision at node 703 without testing. A depth first search would now traverse from node 703 directly to node 706. A breadth first search would first consider nodes 704 and 705 and conclude not to visit them because each is a descendent from a leaf node.

Having selected node 706 at time 830, commander station 10 broadcasts an ID command with BRANCH=001 and MASK=011 at block 216. At time 835 responder station 40 has determined that it is addressed and has begun transmitting response IDR. Simultaneously, responder station 36 determines it is not addressed and remains in state 314. At time 840, shown on FIGS. 13 and 14, commander station 10 has received the response from responder station 40 as a proper message, concluded that only one responder station responded, derived received ARBITRATION NUMBER (101), set BRANCH to received ARBITRATION NUMBER, set MASK to all 1's so that a responder station must match ARBITRATION NUMBER (101) in all bit positions to respond, and begins to perform blocks 232 through 244 in FIG. 10. At time 845, responder station 40 has determined that it is addressed, has decoded a read command, has set its locked-bit 88 in state 326, and has begun generating the read response in state 328. At time 850, commander station 10 has received the response as a proper message. Thus, commander station 10 has conducted a first two-party uninterrupted command-response scenario from time 840 to time 850 with one responder station.

The search by commander station 10 for another responder station proceeds from block 244 to block 224 in FIG. 10. At block 226, another node from FIG. 12 is selected. Having elicited a proper response at node 706, the depth first search proceeds up to the first node having an untested branch, here node 703. Then, down the untested branch to node 707. Having selected node 707 at time 850, commander station 10 broadcasts an ID command with BRANCH=011 and MASK=011 at block 216. At time 855, responder station 36 has determined that it has been addressed and has begun generating an IDR response. At time 860, the response is received by commander station 10 as a proper message. After time 860, events proceed in a manner similar to events from time 840 to time 850, as commander station 10 conducts a second two-party uninterrupted command-response scenario with a second responder station.

At block 224, following the uninterrupted scenario, commander station 10 can conclude that all groups have been tested. On a depth first search, a proper response or no response at a node having BRANCH=MASK indicates all groups have been tested. On a breadth first search, all groups have been tested when an investigation of all levels up to the level having all MASK bits set to '1' yields no node that is not descendent from a leaf node.

In a branch/mask embodiment of the type described above, a responder station concludes that it has been addressed when ARBITRATION NUMBER logically ANDed with MASK is equal to BRANCH. Two other types

of embodiments will now be described that lie within the scope and spirit of the present invention. First, in an example of a high/low embodiment, BRANCH and MASK (as shown in format 142) are replaced with HIGH LIMIT and LOW LIMIT. Each of these limit values has the same precision as the MASK value. Using these limit values, responder station 40 concludes that it is addressed when HIGH LIMIT is greater than or equal to ARBITRATION NUMBER, and LOW LIMIT is less than or equal to ARBITRATION NUMBER. Second, in an example of a limit/bound embodiment, BRANCH and MASK (as shown in format 142) are replaced with a single LIMIT value having the same precision as MASK. Using a value called BOUND which by design choice may be 0 or the maximum permitted by the precision of LIMIT, responder station 40 concludes that it is addressed when ARBITRATION NUMBER is between BOUND and LIMIT, inclusive of both BOUND and LIMIT values.

An example of a limit/bound embodiment is implemented with a structure similar to that already described for the branch/mask embodiment. Subtraction capability or equivalent must be added to ALU 72. Operation of microsequencer 42 must be revised to perform the arithmetic operations outlined above in state 312 shown on FIG. 11. The high/low embodiment is implemented with the structure already described for the limit/bound embodiment.

In FIG. 10 (blocks 210 and 226) commander station 10 specifies a group of responder station addresses. For a branch/mask embodiment, a method using the binary tree of FIG. 12 has already been discussed. For a high/low embodiment, a similar binary tree (not shown) with HIGH and LOW values at each node is used. At the root node, LOW is 0 and HIGH is the maximum value permitted by the precision of the value HIGH. At a node on the lower left from a parent node, the value of LOW is the value of LOW at the parent node and the value of HIGH is a value  $\frac{1}{2}$  the value of HIGH at the parent node discarding a remainder, if any. At a node on the lower right from a parent node, the value of HIGH is the value of HIGH at the parent node and the value of LOW is  $\frac{1}{2}$  the value of HIGH at the parent node plus one. Although a binary tree has been described, a tree having more than two branches at each node can be employed to practice the present invention as is readily apparent to those skilled in the art. Trees with varying number of branches at each node can also be employed. Operation of the high/low embodiment is otherwise identical to operation of the branch/mask embodiment already discussed.

In a limit/bound embodiment, the method used to specify a group of responder station addresses is similar to the method described for a high/low embodiment with a minor variation in the tree. When BOUND is zero, then the value for LOW is not used and the value for HIGH is used as the LIMIT value at each node. When BOUND is a maximum value, then the value for HIGH is not used and the value of LIMIT at each node is the value of LOW. Operation of a limit/bound embodiment is otherwise identical to operation of a branch/mask embodiment already discussed. Note that the command at block 232 on FIG. 10 sets locked-bit 88 to prevent unnecessary collisions when an ID command using LIMIT is broadcast subsequently at block 228.

FIG. 15 is a fibonacci tree diagram for use in an example of an embodiment of the type already described as limit/bound. An advantage of using a fibonacci tree is that the LIMIT value for a node descendent from a parent node can be derived without a multiplication or division operation. In systems where it is desirable to improve calculation speed or

reduce the complexity of circuitry and software at commander station 10, the fibonacci tree is used. An implementation of a high/low embodiment using a fibonacci tree similar to FIG. 15 is within the ordinary skill of the systems design and programming arts.

As described in several embodiments above, a commander station can quickly determine the identity of all responder stations coupled to a common medium at a given time. After the identity of a responder station has been determined, a commander station can conduct uninterrupted communication at any subsequent time using the responder station's ARBITRATION NUMBER. Since the ARBITRATION NUMBER is not unique, there is some risk that at a subsequent time, more than one responder station having a given value for ARBITRATION NUMBER may become coupled to the common medium. For increased accuracy, use of a unique responder station identity, such as the TAG field in format 146 of FIG. 6, may be used for subsequent communication.

When more than one commander station is coupled to a common medium, it is possible for one commander station to thwart the objective of a second commander station. For example, when commander station 10 is attempting to identify all responder stations and commander station 34 issues an IDCG command, commander station 10 may subsequently incorrectly conclude that all responder stations have been identified. Several methods of preventing this incorrect conclusion are available to those skilled in communication and computer programming arts. Exemplary methods include enabling a commander station to monitor commands issued by another commander station to avoid inappropriate conclusions; enabling a commander station 10 to record the TAG fields sent in messages to another commander station and communicate directly with each such responder station, perhaps prior to and so simplifying, the task of identifying all responder stations; modifying the communication protocol used among commander stations; and causing a second commander station to delay its own attempt to identify all responder stations until after a time sufficient for a first commander station to identify all responder stations. The latter suggestion is practical using the media access control scheme of the present invention. It is practical because the time required to identify a worst-case population of responder stations can be predetermined.

The present invention can be implemented in several alternate embodiments. As already discussed, various alternatives are available for common medium 32 including all media that support forms of electromagnetic energy, all sound, vibration, and pressure wave conducting media, and all media capable of transporting variation in chemical concentration, to name a few. If a medium other than radio communication is selected as an embodiment of the present invention, variations in network interface 26 and 60 can be made by those skilled in the arts that apply to the selected medium. Appropriate signal sensors and generators are well known in applications including measurement and control apparatus. Packet synchronization techniques, packet formats, error detection techniques, and error correction techniques may vary or be omitted as a matter of design choice depending on the need for receiver synchronization, the signal to noise properties of the selected media, and the desired simplicity of network interfaces.

Another group of alternative embodiments uses various means to specify a set of responder station identities or designations. The embodiments described above employ an ARBITRATION NUMBER selected from a predetermined range of numbers and expressed as part of a message. For

example, alternate sets of designations include a set of operating modes, a set of modulation techniques, and a range of values used to shift in time all or a portion of a command. Various alternatives are also available for specifying (i.e. addressing) a subset of designations. The branch/mask, high/low, and limit/bound subset addressing techniques can each be applied to one or more parametric quantities related to the above mentioned set designations. For example, if one member of the set is characterized by a bandwidth, a channel frequencies, a phase variation, or a duration in time, then a range of each of these parameters could be described by a branch/mask pair of values.

Various alternatives for transmitting the command signal are within the scope of the present invention. In the embodiments described at length above, the BRANCH and MASK values in the message format characterize the transmitted command signal according to a subset of responder stations to which the command is directed. In addition to the variations in modulation already described, the transmitted signal can be characterized by variation in the spread spectrum chip sequence or initial code within a chip sequence when spread spectrum transmission is employed.

Other characteristics of a command signal can be used to limit or expand the subset of responder stations to which the command is directed. For example, the operation of commands including RD and WR to set locked-bit 88 and the operation of commands including IDG and IDCG to conditionally clear locked-bit 88 show how the command opcode can be used to characterize a command signal according to a selected subset or address range. Alternatively, modulation variations, timing variations, or other message content variations could also be used to set or clear an equivalent of the locked-bit function.

Various means are suitable for use by a responder station to determine whether it is addressed by, i.e. whether it is a member of the subset indicated by, a command signal. Several arithmetic comparison techniques based on message content have been described above. Other means, based on whether the signal received by the responder is received without error, are appropriate when variations in modulation are used to address a subset of responder stations. For example, received signal strength below a threshold over one or more frequencies or at a particular time could cause commands to be received or rejected. Similarly, operation of functions similar to locked-bit 88 as already described and variation in spread spectrum chip sequence could be used to cause commands to be received or rejected.

Within the scope of the present invention, each responder station includes means for establishing a self designation. In the embodiments discussed at length above, the self designation is determined by a random number generator, held in a register, and included in a response packet. Alternative techniques include various means for sampling a random process to acquire an analog parametric value and using either a digital or an analog value to control the functions of network interface 60.

Network interface 60 can be constructed and operated in several alternative embodiments to transmit a response packet in a way characterized by the responder station self designation. All of the following variations could be used in embodiments that fall within the scope of the present invention: variations in the modulation technique, including variation within a range of values used to shift in time all or a portion of a response; variation in the spread spectrum chip sequence or initial code within a chip sequence when spread spectrum transmission is employed; variation in message

content including preamble, postamble, response type indicator e.g. IDR, RDR, and WDR, register contents, status and locked-bit information; and variation based on signal rejection including variation in bandwidth, channel frequencies, signal phase variations, signal duration, or variation in the redundancies used to detect or correct transmission error.

Another group of alternative embodiments uses alternative means for selecting a subgroup in response to collision detection. The tree search method that was described as part of the commander station protocol can be implemented in various ways depending on the selected representation of the tree in commander station memory 18. Binary trees have been described above. Other tree structures including n-ary trees could be employed to perform the commander station identification function in an equivalent manner. Depending on the type of tree selected for representation, the use of strings, arrays, stacks, pointers, linked lists, or derivative memory organizations are feasible and equivalent. Finally, tree search methods include depth first, breadth first, and combinations of both depth and breadth searching.

Each computer used as part of commander station 10 and as part of responder station 40 includes hardware and software designed to conduct the protocol shown and described in FIGS. 10 and 11 respectively. Variations in the extent and complexity of hardware and software are well known by designers of ordinary skill in communication and computer arts. Equivalent hardware can include the general purpose computer such as an IBM PC; a calculator, such as an HP21C; the special purpose computer, such as application specific automated controllers used in weighing systems; the microprocessor based system, such as a circuit using an Intel 8048; the microsequencer based system using programmable devices and logic devices; and the integrated circuit or chip set, such as developed from a cell library using semiconductor device design techniques. Variations in the extent and complexity of software compatible with one or more of the above mentioned hardware variations are also well known by the programmer of ordinary skill.

The systems designer of ordinary skill chooses to implement each control function in either hardware or software or a combination of both. A computer is said to conclude a certain result when it has determined the state of a control function. When a control function is implemented using a computer system, variations in the form of the result of the control function are well known. For example, a parameter that results from a first control function and is relied upon by a second control function can take the form of a signal when the second control function is in part hardware or the form of a pointer, value, or symbol stored in a register or memory when the second control function is in part software.

The present invention has been described in the preferred embodiment. Several variations and modifications have also been described and suggested. Other embodiments, variations, and modifications known to those skilled in the art may be implemented without departing from the scope and spirit of the invention as recited in the claims below.

What is claimed is:

1. For communication between a commander communication transceiver and a plurality of responder communication transceivers, a method for each responder transceiver to establish an address number so that it can be individually addressed by the commander transceiver, comprising the steps of:

- a) establishing a set of possible addresses for responder transceivers;
- b) the commander transceiver transmitting a generate message requesting responder transceivers who receive the message to select a new address; and

- c) upon receipt of the generate message, each responder transceiver choosing a random address from the set of possible addresses, wherein each responder transceiver chooses its address independently of addresses chosen by all other responder transceivers;
- d) whereby the addresses chosen by the responder transceivers are randomly distributed within the set of possible addresses, and whereby the commander transceiver controls when the responder transceivers choose new addresses.
2. A method according to claim 1, further comprising the subsequent steps of:
- each responder transceiver transmitting an identification message containing its chosen address;
- upon receipt of identification messages from different responder transceivers containing the same chosen address, the commander transceiver transmitting an address reset message; and
- upon receipt of the address reset message, each responder transceiver choosing a new address.
3. A method according to claim 1, further comprising the subsequent steps of:
- the commander transceiver transmitting an identification request message requesting responder transceivers having addresses within a specified first subset of addresses to identify themselves without selecting a new address, the first subset being less than the entire set of possible addresses; and
- upon receipt of the identification request message, each responder transceiver
- i) determining whether its previously chosen address is within the first subset, and if it is, then
  - ii) transmitting an identification message containing its chosen address.
4. A method according to claim 3, further comprising the subsequent steps of:
- upon receipt of identification messages from different responder transceivers containing the same chosen address, the commander transceiver transmitting an address reset message addressed to a specified second subset of addresses, wherein the second subset includes the first subset, the second subset being less than or equal to the entire set; and
- upon receipt of the address reset message, each responder transceiver
- i) determining whether its most recently chosen address is within the second subset, and if so, then
  - ii) choosing a new address.
5. A method according to claim 1, wherein:
- the generate message specifies for response a first subset of addresses within the set of possible addresses, said first subset being less than or equal to the entire set of possible addresses; and
- after the step of choosing a random address, each responder transceiver
- i) determining whether its chosen address is within the first subset specified in the received first message, and if so, then
  - ii) transmitting an identification message containing its chosen address.
6. A method according to claim 5, further comprising the subsequent steps of:
- the commander transceiver transmitting an identification request message requesting responder transceivers having addresses within a specified second subset of

- addresses to identify themselves without selecting a new address, the second subset being less than the entire set; and
- upon receipt of the an identification request message, each responder transceiver
- i) determining whether its previously chosen address is within the second subset, and if it is, then
  - ii) transmitting an identification message containing its chosen address.
7. A method according to claim 5, further comprising the subsequent steps of:
- upon receipt of identification messages from different responder transceivers containing the same chosen address, the commander transceiver transmitting an address reset message; and
- upon receipt of the address reset message, each responder transceiver choosing a new address.
8. A method according to claim 5, further comprising the subsequent steps of:
- upon receipt of identification messages from different responder transceivers containing the same chosen address, the commander transceiver transmitting an address reset message addressed to a specified third subset of addresses, wherein the third subset includes the first subset, the third subset being less than or equal to the entire set; and
- upon receipt of the address reset message, each responder transceiver
- i) determining whether its previously chosen address is within the third subset, and if it is, then
  - ii) choosing a new address.
9. A method of addressing messages from a commander station communication transceiver to a selected one or more of a plurality of responder station communication transceivers, comprising the steps of:
- a) establishing a first predetermined number of bits to be used for unique identification numbers, and establishing for each responder transceiver a unique identification number having said predetermined number of bits;
  - b) establishing a second predetermined number of bits to be used for arbitration addresses, the second number being less than the first number, and establishing a set of possible arbitration addresses having said second number of bits;
  - c) the commander transceiver transmitting a generate message requesting responder transceivers who receive the generate message to select a new arbitration address;
  - d) the commander transceiver transmitting an identification request message requesting responder transceivers to respond identifying themselves, wherein the message specifies for response a first subset of arbitration addresses within the set of possible arbitration addresses, said first subset being less than the entire set; and
  - e) after receiving the generate message, each responder transceiver choosing an arbitration address having the second number of bits, wherein each responder transceiver chooses its arbitration address independently of arbitration addresses chosen by all other responder transceivers; and
  - f) after receiving the identification request message, each responder transceiver
- i) determining whether its chosen arbitration address is within the first subset specified in the identification request message, and if it is, then



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ii) transmitting an arbitration identification message containing its chosen arbitration address and a unique identification message containing its unique identification number.

10. A method according to claim 9, wherein

after receiving an arbitration identification message containing the arbitration address of a responder transceiver, the commander transceiver transmitting a command message containing that arbitration address; and the step of each responder transceiver transmitting the unique identification message is delayed until the responder transceiver receives said command message from the commander transceiver.

11. A method according to claim 9, wherein the unique identification message transmitted by each responder transceiver includes both that transceiver's arbitration address and its unique identification number.

12. A method according to claim 9, further comprising the subsequent steps of:

upon receipt of arbitration identification messages from different responder transceivers containing the same arbitration address, the commander transceiver transmitting an address reset message; and

upon receipt of the address reset message, each responder transceiver choosing a new arbitration address.

13. A method according to claim 9, further comprising the subsequent steps of:

upon receipt of identification messages from different responder transceivers containing the same chosen address, the commander transceiver transmitting an address reset message addressed to a specified second subset of addresses, wherein the second subset includes the first subset, the second subset being less than or equal to the entire set; and

upon receipt of the address reset message, each responder transceiver

- i) determining whether its most recently chosen address is within the second subset, and if so, then
- ii) choosing a new address.

14. A method of establishing an address number in a responder station communication transceiver so that it can be individually addressed by a commander station communication transceiver, comprising the steps of:

- a) establishing a set of possible responder transceiver addresses;
- b) the commander transceiver transmitting a first command message;
- c) after receipt of the first command message, each responder transceiver
  - i) choosing an address independently of addresses chosen by all other responder transceivers, and then
  - ii) transmitting an identification message containing its chosen address;
- d) the commander transceiver
  - i) detecting whether it receives identification messages from different responder transceivers containing the same address, and if so, then
  - ii) transmitting a second command message; and
- e) after receipt of the second command message, each responder transceiver choosing a new address;
- f) whereby the commander transceiver can individually address a message to any of said responder transceivers by using the responder transceiver's chosen address.

15. A method according to claim 14, wherein each step of choosing an address comprises choosing a random address within the set of all possible addresses.

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16. A method of establishing an address number in a responder station communication transceiver so that it can be individually addressed by a commander station communication transceiver, comprising the steps of:

- a) establishing a set of possible responder transceiver addresses;
- b) the commander transceiver transmitting a first command message requesting responder transceivers who receive the message to select a new address and to respond identifying themselves, wherein the message specifies for response a subset of addresses within the set of possible addresses, the subset being less than or equal to the entire set;
- c) after receipt of the first command message, each responder transceiver
  - i) choosing an address independently of addresses chosen by all other responder transceivers, and then
  - ii) if the chosen address is within the subset specified in the first command message, transmitting an identification message containing the chosen address;
- d) the commander transceiver
  - i) detecting whether it receives identification messages from different responder transceivers containing the same address, and if so, then
  - ii) transmitting a second command message addressed to a second subset of addresses, wherein the second subset includes the first subset, the second subset being less than or equal to the entire set of possible addresses; and
- e) after receipt of the second command message, each responder transceiver
  - i) determining whether its previously chosen address is within the second subset, and if it is, then
  - ii) choosing a new address
- f) whereby the commander transceiver can individually address a message to any of said responder transceivers by using the responder transceiver's chosen address.

17. A method according to claim 16, wherein each step of choosing an address comprises choosing a random address within the set of all possible addresses.

18. A method for a commander communication transceiver to poll a plurality of responder communication transceivers to determine their addresses, comprising the steps of:

- (A) establishing a set of possible addresses for the responder transceivers;
- (B) the commander transceiver transmitting a Clear Polling Lock message;
- (C) upon receipt of the Clear Polling Lock message, each responder transceiver clearing a Polling Lock flag in a memory device within that responder transceiver;
- (D) the commander transceiver transmitting a first polling message requesting responder transceivers having addresses within a specified first subset of addresses to identify themselves, the first subset being less than or equal to the entire set of possible addresses;
- (E) upon receipt of the first polling message, each responder transceiver
  - 1) determining whether its address is within the first subset, and if it is, then
  - 2) transmitting an identification message containing its address;
- (F) the commander transceiver determining whether it receives an identification message from a responder transceiver which is not corrupted by a collision with another message, and if so, then

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- 1) transmitting a Set Polling Lock message addressed to the responder transceiver from whom the commander transceiver received the uncorrupted identification message, and
  - 2) subsequently transmitting a second polling message 5 requesting responder transceivers having addresses within a specified second subset of addresses to identify themselves, wherein the second subset includes the first subset, the second subset being less than or equal to the entire set of possible addresses; 10
- (G) each responder transceiver, upon receipt of a Set Polling Lock message addressed to that responder transceiver, setting the Polling Lock flag within that responder transceiver;
- (H) upon receipt of the second polling message, each responder transceiver 15
- 1) determining whether its address is within the second subset and whether its Polling Lock flag is clear, and if both conditions are true, then
  - 2) transmitting an identification message containing its address; and 20
- (I) the commander transceiver

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- 1) determining whether it receives, in response to the second polling message, a plurality of identification messages which collide with each other so as to be unintelligible to the commander transceiver, and if so, then
  - 2) repeating steps (F) through (I), using in each repetition a progressively smaller second subset, until the commander transceiver does not receive unintelligible, colliding identification messages in response to the second polling message.
19. A method according to claim 18, further including the steps of
- the commander transceiver determining whether any of the identification messages transmitted by the responder transceivers contain duplicate addresses, and if so, then
- i) the commander transceiver transmitting a Clear Polling Lock message, and
  - ii) the commander transceiver transmitting a message requesting the responder transceivers to choose new addresses.

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